Identifying Slow Queries, and Fixing Them

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Introduction

- Stephen Frost
 - System Architect/Designer
 - DBA, Unix Administrator
 - PostgreSQL/PostGIS Hacker
 - Added Roles in 8.1, Column-level Privs in 8.4
- Noblis, Inc.
 - Nonprofit science, technology and strategy organization
 - http://www.noblis.org



Finding the slow ones

(Queries...)



Monitor your systems!

PostgreSQL Logs

- Configure what gets logged!
- Log checkpoints, connections, DDL statements! Perhaps more..
- Your favorite monitoring solution
 - Availability, Alarm-based (eg: Nagios, w/ check_posgres)
 - Performance measuring (eg: munin, stats w/ pg_bouncer)
 - PgFouine for log file analysis
- check_postgres script
 - Includes lots of valuable checks
 - Bloat checking
 - Idle connection warnings
 - Number of WAL files (in case archiving fails)
 - Can integrate w/ munin/cacti/MRTG too!



Finding Slow Queries

- postgresql.conf
 - log_min_duration_statement just needs reload
 - Lots of other logging options:
 - log_line_prefix
 - log_connections / log_disconnections
 - log_duration
 - log_lock_waits
 - log_statement
 - track_functions

Reviewing PG logs

- LOG: duration: 448.495 ms statement: select generate_series(1,100000);
- What's in that duration?
- Difference with psql's \timing option



PG duration logging

- More PG logs
 - Just log_min_duration_statement:
 - LOG: duration: 448.495 ms statement: select generate_series(1,100000);
 - vs. log_statement = all && log_min_duration_statement:
 - LOG: statement: select generate_series(1,100000);
 - LOG: duration: 513.041 ms
 - vs. log_statement = none && log_min_duration_statement && log_duration:
 - LOG: duration: 0.659 ms
 - LOG: duration: 457.366 ms statement: select generate_series(1,1000000);
 - If you can afford log_statements=all and log_duration you can gather lots of info, but it's not free to log at that level (typically not done in high-transaction production systems)
 - log_min_duration_statement gives 'best of both worlds'- just log the slow ones, but be careful what other options you have enabled or it may get confusing
 - Lots of fast queries, done sequentially, can also make things (page loads) slow!



Now we've found them ...

Why are they slow?



Understanding why queries are slow

- The "easy" stuff-
 - Poor PG configuration
 - Dead tuples / bloat
- The next level- Database Magic



Poor PG Configuration (you used the defaults...)

- Important PG GUCs (configuration options):
 - work_mem
 - maintenance_work_mem
 - effective_cache_size
 - shared_buffers
 - checkpoint_segments
- Watch for differences between Prod & Dev
 - Need to understand them, if any
 - May get different plans if different
 - "Unseen" differences
 - Statistics data may be different
 - Different hardware
 - Warm-up Time



Dead Tuples / Bloat

- VACUUM marks records reusable, if possible
 - Reusable tuples will be used for new inserts, etc
 - However, PG has to handle those tuples on queries
- Records marked as deleted but not reusable yet

 Ongoing transactions
- Bloat can exist in both tables and indexes
- check_postgres.pl
 - Can identify bloat in tables/indexes
 - Some bloat is GOOD, but too much will make queries slower (lots of extra/unnecessary data to process)
- CLUSTER will re-write a table and eliminate dead tuples.



Database Magic, or how it works

- There is no magic here, sadly.
- Getting data:
 - Sequentially step through EVERY record
 - SeqScan Node
 - Bulk, very fast at going through a table
 - Pick out SPECIFIC records, using an index
 - Index Scan Node
 - Very slow for bulk data
 - Can return data in-order
 - Index needs to be there..



More Magic

- Putting things together (Joins)
 - Loop through and scan table for match
 - Nested Loop Node
 - Works for **small** data sets
 - Order two tables, then walk through each Merging them
 - Merge Join Node
 - Requires sorted inputs
 - Good for bulk operations, esp. work loads that won't fit in memory
 - Build a hash table (of the smaller table) then step through
 - Hash Join Node
 - Requires lots of memory
 - Very fast, but slow to start
- Adding it all up (Aggregates)
 - Look at all rows that qualify
 - Can be very expensive



What's the best plan?

- It depends!
- How's the database know?
 - Gathers statistics using ANALYZE
 - Automatically done by auto-vacuum
- What if the database (aka- the stats) are wrong?
 - You get bad plans!
 - Look for differences in row estimates from explain analyze:
 - Index Scan using my_idx on my_table (cost=0.00..5719.56 rows=9055 width=10) (actual time=0.015..87.689 rows=163491 loops=1)
 - May need to adjust statistics target



What plan did PG decide to use?!



Understanding "explain"

• explain output:

000	Default	\Box
Default		
postgres=# explain select	* from pg_class a join pg_namespace b on (a.relnamespace = b.oid); QUERY PLAN	
-> Hash (cost=1.061	bace = b.oid) ss a (cost=0.0010.81 rows=281 width=194)	

- Node types: Hash, Hash Join, Seq Scan
- Lots of other node types
- What is the cost?



More Explain

• With "where"

00	Default	
Default		
postgres=# explain select	* from pg_class a join pg_namespace b on (a.relnamespace = b.oid) where relname = 'pg_class'; QUERY PLAN	Ô
Index Cond: (rein	espace = b.oid) g_class_relname_nsp_index on pg_class a (cost=0.008.27 rows=1 width=194) name = 'pg_class'::name)	
-> Seq Scan on pg_name (5 rows)	space b (cost=0.001.06 rows=6 width=117)	0
Time: 1.221 ms postares=#		•

- Very different plan!
- More nodes: Nested Loop, Index Scan
- Lower cost, much fewer rows



Understanding "explain analyze"

• explain analyze output

0 0	Default	\Box
Default		
postgres=# explain analyzo	e select * from pg_class a join pg_namespace b on (a.relnamespace = b.oid); QUERY PLAN	<u> </u>
Hash Cond: (a.relnames) -> Seq Scan on pg_clas -> Hash (cost=1.06 Buckets: 1024 Bo	5.81 rows=281 width=307) (actual time=0.0741.960 rows=281 loops=1) pace = b.oid) ss a (cost=0.0010.81 rows=281 width=194) (actual time=0.0120.583 rows=281 loops=1) 1.06 rows=6 width=117) (actual time=0.0410.041 rows=6 loops=1) atches: 1 Memory Usage: 1kB pg_namespace b (cost=0.001.06 rows=6 width=117) (actual time=0.0050.018 rows=6 loops=1)	
(7 rows) Time: 3.617 ms postgres=#		↓ ↓ //.

Lots more info- actual times, per-node info, memory usage!

Two times? backand runtime neal timing



More explain analyze

• Explain analyze with where output

0 0	Default	\Box
Default		
postgres=# explain analyze	: select * from pg_class a join pg_namespace b on (a.relnamespace = b.oid) where relname = 'pg_class'; QUERY PLAN	ĥ
Nested Loop (cost=0.00 Join Filter: (a.relname	9.40 rows=1 width=307) (actual time=0.0520.072 rows=1 loops=1) space = b.oid)	
Index Cond: (rein	j_class_relname_nsp_index on pg_class a (cost=0.008.27 rows=1 width=194) (actual time=0.0220.026 rows=1 loops=1 ame = 'pg_class'::name))
	space b (cost=0.001.06 rows=6 width=117) (actual time=0.0060.021 rows=6 loops=1)	
Total runtime: 0.147 ms (6 rows)		Q
Time: 1.280 ms postgres=#		▼

- Back to the other plan, with actuals, total runtime
- Still a seqscan on pg_namespace..?



Other explain output options

- Other output options
 XML, JSON, YAML
- Tools to analyze explain output
 - PgAdmin3
 - explain.depesz.com
- PG Log file analyzer, includes tracking timing info – pgFouine



Automating collection of "explain"s

- auto_explain
 - Logs explain info for long queries
- Enabling:
 - shared_preload_libraries = 'auto_explain'
 - set explain.log_min_duration = 50;
 - Can also output 'explain analyze' (expensive!)
 - Set explain.log_nested_statements = true;
 - Considers logging plans which are inside functions



But how to change the plan? How to fix the queries..?!



Fixing Them

- Low-hanging fruit, but catch a lot..
 - Query returning 1 row using SeqScan? Check for an index
 - MergeJoin used for "small" data sets? Check work_mem
 - Nested Loop used with large set? Bad row estimates?
 - Make sure analyze is being done
 - Increase statistics target for the relations if possible
 - DELETE's slow? Make sure you have indexes on Foreign Keys
- Harder items:
 - Check over your long-running queries
 - Use stored procedures / triggers
 - Partitioning large tables
 - Consider Partial Indexes / Functional Indexes



SeqScan returns 1 row

• Lack of index on username

0 0			Default	\Box
<u>Å</u>	Default	Default		
postgres=#	explain analyze		where username = 'user2'; RY PLAN	1
Filter:		cost=0.00738.00 rows=1 w ext = 'user2'::text)	vidth=268) (actual time=0.01911.916 rows=1 loops=1)	
postgres=# count	select count(*)) from customers;		
20000 (1 row)				
postgres=#	1			



Using an index scan

• Much better performing, what about like?

00		Default	\Box
Default	Default		
postgres=# explain analyze	select * from customers	where username = 'user2'; QUERY PLAN	î
Index Scan using ix_cust_ Index Cond: ((username) Total runtime: 0.090 ms (3 rows) postgres=#		cost=0.008.27 rows=1 width=268) (actual time=0.0330.038 rows=1 loop	os=1)
00		Default	\Box
Default	Default		
postgres=# explain analyze		where username like 'user212%'; RY PLAN	ſ
Filter: ((username)::te Total runtime: 12.756 ms		width=268) (actual time=0.15512.688 rows=11 loops=1)	
(3 rows)			



Text Pattern Searches

• Need an appropriate index

000		Default	\bigcirc
Default	Default		
postgres=# create index cu CREATE INDEX	st_text_ops_idx on custom	ers (username text_pattern_ops);	ŕ
		where username like 'user212%'; where username like 'user212%'; QUERY PLAN	
Filter: ((username)::te -> Bitmap Index Scan c	ext ~~ 'user212%'::text) on cust_text_ops_idx (cos	rows=2 width=268) (actual time=0.0450.073 rows=11 loops=1) st=0.004.28 rows=3 width=0) (actual time=0.0280.028 rows=11 loops=1) 212'::text) AND ((username)::text ~<~ 'user213'::text))	

- Pattern needs to be anchored and simple
- PG has excellent Full Text Search & Indexing



MergeJoin for 'small' data

• Merge vs Hash and work_mem

For the best of reasons

000		Default)
Default	Default			Ī
postgres=# explain analyze s	elect * from orders join cust	comers using (customerid); QUERY PLAN		Ĩ
Merge Cond: (orders.custo -> Index Scan using ix_o	omerid = customers.customerid) order_custid on orders (cost=	=0.00720.24 rows=12000 width=30) (ad	loops=1) ctual time=0.01331.185 rows=12000 loops=1)) (actual time=0.01053.652 rows=23004 loops=1)	
postgres=# 🛛				1
000		Default	C)
Default	Default			
postgres=# explain analyz	e select * from orders joi	n customers using (customerid); QUERY PLAN		ĥ
Hash Cond: (customers. -> Seq Scan on custom -> Hash (cost=220.00 Buckets: 2048 B	customerid = orders.custom ers (cost=0.00688.00 ro 220.00 rows=12000 width= atches: 1 Memory Usage: 7 orders (cost=0.00220.00	ws=20000 width=268) (actual time= 30) (actual time=53.49753.497 r %1kB	0.00839.573 rows=20000 loops=1)	
postgres=# 🏾			-	1
noblís				

Nest Loops can be good

• For small sets

000	Default	\supset
Default	Default	
postgres=# explain analyza	select * from orders join customers using (customerid) where username = 'user21'; QUERY PLAN	ĥ
-> Index Scan using co Index Cond: ((uso -> Index Scan using i)	16.55 rows=1 width=294) (actual time=0.0310.043 rows=1 loops=1) st_text_ops_idx on customers (cost=0.008.27 rows=1 width=268) (actual time=0.0130.016 rows=1 loops=1) rname)::text = 'user21'::text) _order_custid on orders (cost=0.008.27 rows=1 width=30) (actual time=0.0060.010 rows=1 loops=1) omerid = customers.customerid)	
postgres=# set enable_nest SET postgres=# explain analyze	loop to false; select * from orders join customers using (customerid) where username = 'user21'; QUERY PLAN	
Hash Cond: (orders.cust -> Seq Scan on orders -> Hash (cost=8.27t Buckets: 1024 Bo -> Index Scan us	3.29 rows=1 width=294) (actual time=39.76046.393 rows=1 loops=1) omerid = customers.customerid) (cost=0.00220.00 rows=12000 width=30) (actual time=0.01023.004 rows=12000 loops=1) .27 rows=1 width=268) (actual time=0.0260.026 rows=1 loops=1) tches: 1 Memory Usage: 1kB ing cust_text_ops_idx on customers (cost=0.008.27 rows=1 width=268) (actual time=0.0120.015 rows=1 loops=1) ((username)::text = 'user21'::text)	



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Slow DELETE

For the best of reasons

• Explain analyze on delete, what's the difference?

000		Default		\bigcirc
Default	Default			
postgres=# explain analyze	e delete from customers wi	here username = 'user1'; QUERY PLAN		
-> Index Scan using cu	ust_text_ops_idx on custon rname)::text = 'user1'::* _cust_hist_customerid: t	text) ime=0.349 calls=1	=0 loops=1) =6) (actual time=0.0390.043 rows=1 loo	ps=1)
000		Default		\ominus
Default	Default			
postgres=# explain analyze	: delete from customers wl	here username = 'user1'; QUERY PLAN		
-> Index Scan using cu	ust_text_ops_idx on custon rname)::text = 'user1'::* _cust_hist_customerid: t	text) ime=0.064 calls=1	=0 loops=1) =6) (actual time=0.0330.037 rows=1 loo	ps=1)
			1	

Prepared queries

- They're good, honest
- Plan once, run many
 - Not as much info to plan with, plans may be more stable
 - Variables aren't substituted in until execution
 - No constraint exclusion though
- How to explain/explain analyze:
 - prepare q as select * from table where x =\$1;
 - explain execute q('myid');
 - explain analyze execute q('myid');
- Placeholders in explain output (\$1 instead of 'myid')



Query Review

- select count(*) from table;
 - Expensive, must check every record in the table
- select * from table;
 - Returns every row, do you really need them all?
 - Order By / Limit can help PG optimize queries!
- select * from table where id = 1;
 - Do you need every column? Wide columns cost / TOAST
- select * from a, b, c where a.x = b.x;
 - Missing join condition for c!
 - Cartesian product with a/b to c
 - Use join syntax:
 - select * from a join b using (x) join c using (x);



More Queries

- select * from x where myid in (select myid from big_table);
 - Turn it into a join:
 - select x.* from x join big_table using (myid);
- select * from x where myid not in (select myid from big_table);
 - Left-join instead:
 - select x.* from x left join big_table using (myid) where big_table.myid is NULL;
 - Not exists also:
 - select * from x where not exists

(select * from big_table where big_table.myid = x.myid)



More queries...

- Expensive to generate table? Use CTE (Common Table Expressions, aka WITH)
- select *,

(select sum(my_expensive_view.x) from my_expensive_view)
from my_expensive_view;

- WITH my_view AS (select * from my_expensive_view), my_sums AS (select sum(my_view.x)) select my_view.*, my_sums.sum from my_view, my_sums;
- CTEs can also be used to implement recursion!



Really need fast count(*)?

- Does it have to be accurate or just an estimation?
 Look at pg_class.reltuples for an estimate
- Use a trigger if it needs to be accurate
 - Handle bulk-loading independently though
- It's a trade-off
 - Faster to get count(*) information
 - Slower to insert/update the table
- create function my_count_func() returns trigger as \$_\$

BEGIN

```
UPDATE my_count = my_count + 1;
```

```
RETURN NEW;
```

END \$_\$ LANGUAGE 'plpgsql';

create trigger my_count_trig after insert on my_table for each row execute procedure my_count_func();



What else can be done?

- Tuning PG GUCs
 - work_mem default 1MB is wayyy small
 - maintenance_work_mem default 16MB small
 - effective_cache_size default 128MB
 - You have a server with 256MB of memory..?
 - shared_buffers default 24MB
 - This is a real killer.. bump it to 1-2G, at least, on a server w/ >4G RAM, up to 8GB (don't go above that w/o good testing..).
- Partial Indexes / Functional Indexes
- Improving statistics / analyze / auto-vacuum
- Tuning the background writer
 - Consider making it more aggressive for heavy write loads
- Invest in hardware
 - Lots of memory (adjust shared_buffers..)
 - SSDs / Battery-Backed Write Cache RAID



Questions?

